

Automated Testing of Stateful Network Protocol Implementations in the IoT

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Background: IoT nodes exchange data via network protocols

- Protocol implementations often contain software bugs
- Some of these bugs (e.g. buffer overflows) are exploitable
- Problematic since IoT operating systems have few exploit mitigations

Goal: Automatically find such bugs in network modules

 \Rightarrow Emerging method for this purpose: symbolic execution





Idea: Enumerate reachable paths based on specific input source

- SW executed with symbolic values, represent set of concrete values
- Symbolic values are continuously constrained during execution
- Constraints on current path: path constraints (PC)





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DSE: Dynamic Symbolic Execution

- Concrete execution drives symbolic execution
- Track symbolic constraints alongside concrete execution
- Branches are collected; later negated with an SMT solver



Dynamic Symbolic Execution (1/3)



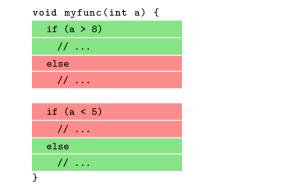
Example:

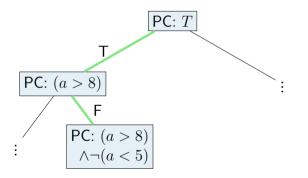
```
void myfunc(int a) {
    if (a > 8)
        // ...
    else
        // ...
    if (a < 5)
        // ...
    else
        // ...
}</pre>
```





Example: Execution trace for myfunc with input a = 9



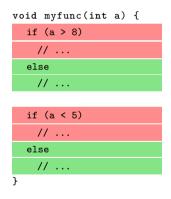


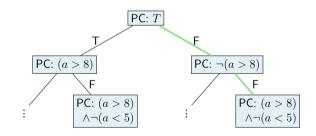




Exploration: Negate unexplored branch $\neg(a > 8)$, solve resulting query

 \Rightarrow Restart execution with concrete input (e.g. a = 8)



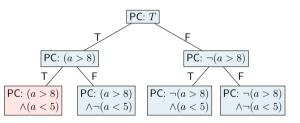




Goal: Ideally discover all execution paths

 \Rightarrow Repeat until all branches have been negated

```
void myfunc(int a) {
    if (a > 8)
        // ...
    else
        // ...
    if (a < 5)
        // ...
    else
        // ...
}</pre>
```



















MQTT-SN: Stateful protocol for data exchange in the IoT

- Certain code can only be tested by establishing a state first
- For example, subscribing to a specific topic
- Results in a large state space for symbolic execution
 - \Rightarrow Cannot be fully explored using symbolic execution





Goal: Discovering "interesting" execution paths first ⇒ Observation: Many inputs are rejected early on

Approach: Partially specify protocol message format

- Embedded domain specific language (EDSL)
- Based on the Scheme programming language

```
(define-input-format (suback id)
 (make-uint 'len 8 8)
 (make-uint 'type 8 MQTT-SUBACK)
 (make-symbolic 'flags 8)
 (make-symbolic 'topicid 16)
 (make-uint 'msgid 16 id)
 (make-symbolic 'code 8
   `((And
        (Uge ,code 0)
        (Ule ,code 3)))))
```

Figure: Message format for MQTT-SN SUBACK.





Challenge: MQTT-SN is a stateful protocol \Rightarrow Message format depends on protocol state

Approach: Also describe protocol state machine

- With a separate Scheme-based EDSL
- Advance protocol state based on received messages
- Return new symbolic message depending on state

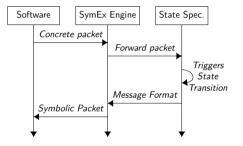


Figure: Overview of message format exchange.



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Needed: MQTT-SN state machine description

- Described as a finite-state machine
- Transitions based on input packet
- Each transitions returns a response format

```
(define-state-machine mott-machine
     (start pre-connected)
3
     (define-state (pre-connected input)
Λ
5
       . . . )
7
     (define-state (connected input)
       (switch (mqtt-msq-type input)
8
         ((SUBSCRIBE)
9
          (-> (make-resp (suback-fmt (msg-id input)))
              subscribed))
         ((DISCONNECT)
12
          (-> (make-resp disconn-fmt)
13
              disconnected))
14
         . . . ) )
15
16
18
```

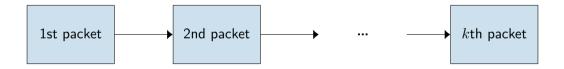
Figure: Excerpt of the MQTT-SN state specification.





Problem: Need to reason about sequence of packets

 \Rightarrow Further increase of state space

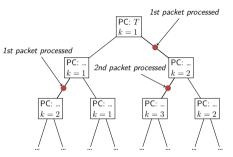






Simplified Algorithm:

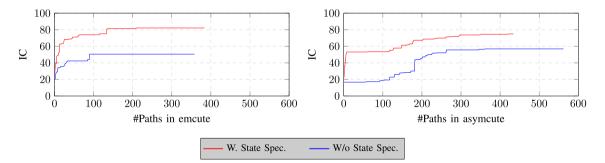
- 1. Explore program up to a sequence length of \boldsymbol{k}
- 2. Restart execution when packet k was processed
- 3. When coverage is stagnant: Increment k
 - \Rightarrow Partially explored paths are re-executed continuously







Research Question: Does our symbolic execution approach improve coverage? \Rightarrow Experiments with RIOT's MQTT-SN implementations

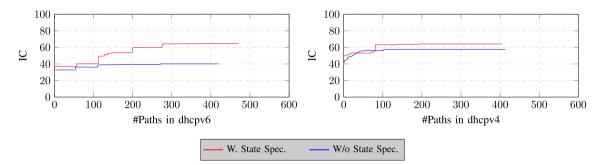






Research Question: Is the approach applicable to other protocols?

 \Rightarrow Experiments with RIOT's and Zephyr's DHCP implementations







Bugs Found:

- 1. #18307: out-of-bounds read in dhcpv6 module
- 2. #18289: missing mutex_unlock in asymcute
- 3. #18434: null pointer dereference in asymcute

Future Work:

- Integrate protocol rules into specification?
- Assessment of created protocol specifications



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Key Insight: High coverage in complex network modules via symbolic execution

 \Rightarrow With comparatively little manual effort

Contributions:

- 1. Input specification language for message formats $^{1} \$
- 2. Specification language for protocol state machines²
- 3. Enhanced version of SymEx-VP with new exploration engine³

More Information: Sören Tempel, Vladimir Herdt, and Rolf Drechsler. *Specification-based Symbolic Execution for Stateful Network Protocol Implementations in the IoT*. IEEE Internet of Things Journal, 2023.

¹https://github.com/agra-uni-bremen/sisl

²https://github.com/agra-uni-bremen/sps

³https://github.com/agra-uni-bremen/sps-vp